

EMBEDDED SYSTEMS REVERSE ENGINEERING

Week 2

Debugging and Hacking a Live System (Pico 2)




Connecting via OpenOCD & GDB, Understanding Memory Constraints (Flash vs. SRAM), Register Manipulation & Hijacking, Writing to Live Memory

What is Live Hacking?

Definition: Modifying a program's behavior *while* it is executing on real hardware, without recompiling the source code.



Why do we do it?

-  **Security Research:** Testing for vulnerabilities in real-time.
-  **Malware Analysis:** Watching how malicious code behaves when it thinks no one is looking.
-  **Debugging:** Fixing bugs in systems that cannot be easily restarted or reprogrammed.

The Target Program & Goal



The Code: 0x0001_hello-world.c



Initializes I/O (stdio_init_all).



Enters an infinite loop.



Prints "hello, world\r\n" forever.

```
int main(void) {  
    stdio_init_all();  
    while (true)  
        printf("hello, world\r\n");  
}
```



The Mission



Intercept the program mid-execution.



Change the output string to "hacky, world".



Constraint: We cannot modify the source code or recompile. We must manipulate the binary state directly.

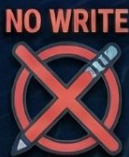
Memory Layout Review (Critical for Hacking)

To hack successfully, we must know *where* we can write data.



Flash Memory (XIP)

- **Address:** Starts at 0x10000000
- **Contents:** Program code and constant strings ("hello, world").
- **Property: READ-ONLY** during runtime. We cannot overwrite the string here.



SRAM (Static RAM)

- **Address:** Starts at 0x20000000
- **Contents:** Stack, Heap, Variables.
- **Property: READ-WRITE.** This is our playground for injecting new data.

Constraint: We cannot



Flash (XIP) - 0x10000000 - Read-Only

SRAM - 0x20000000 - Read-Write

← Data Injection Playground →

The Attack Plan

How we will hijack the program execution:

1. Connect



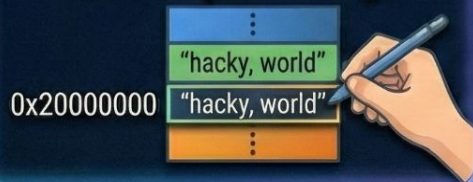
2. Halt



3. Trap



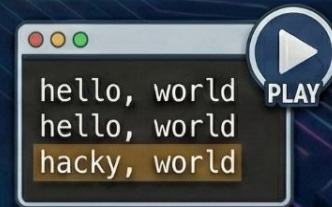
SRAM (Writable)



4. Inject



5. Hijack



6. Resume

Tool Theory: The Debug Chain



1. OpenOCD (Open On-Chip Debugger)

- The bridge between your PC and the Pico 2 hardware.
- Translates high-level commands into electrical signals (SWD/JTAG).



PC



OpenOCD
Bridge



2. GDB (GNU Debugger)

- The “Brain.” Connects to OpenOCD.
- Allows us to stop execution, read memory, and modify registers.



GDB



Memory



Registers



Execution
Control



Pico 2



3. Serial Monitor (PuTTY)

- The “Eyes.” Shows us the actual output coming from the Pico 2.



PuTTY

Analyzing the Target (Disassembly)



GDB Disassembly Output (x/5i 0x10000234)

```
(gdb) x/5i 0x10000234
0x10000234 <main>:      push    {r3, lr}
0x10000236 <main+2>:    bl      0x1000156c <stdio_init_all>
0x1000023a <main+6>:    ldr     r0, [pc, #8]
0x1000023c <main+8>:    bl      0x100015fc <__wrap_puts>
0x10000240 <main+12>:   b.n     0x1000023a <main+6>
```

Loads address of
"hello, world" into r0.

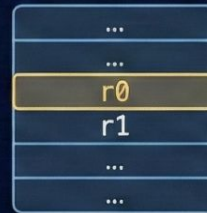
Calls the print
function (puts).



The Critical Vulnerability

The function **puts** takes one argument (the string to print). In ARM architecture, the **first argument is always passed in register r0**.

CPU register



String Address
(e.g., 0x100019cc)

"hello, world\r"

puts(r0)



If we change r0 at address **0x1000023c**, we change what prints.



The Failed Attempt (Why `malloc` is missing)



The Mistake & Error

```
(gdb) set $r0 = "hacky, world\r"  
evaluation of this expression requires the program to have a function "malloc".
```

Failed!

Why it fails:



Program
`malloc()`

Standard Computers (OS)

Operating
System
(Memory
Manager)



RAM

Standard computers have an **Operating System (OS)** to handle memory allocation (`malloc`).

Microcontrollers (Bare Metal)

Even though the Pico C SDK includes `malloc()` and a complete C standard library (newlib/picolibc), GDB cannot use it directly. GDB's automatic function evaluation fails for two reasons:



GDB
`malloc()`



SRAM



(Compiler/
Linker)

1. Linker Optimization (Dead Code Elimination)

The function `malloc()` is removed from the final binary if never explicitly called in your code to save flash memory.



HALTED

2. Halted Processor State

To execute `malloc()`, the processor must run complex code. GDB operates when the core is halted, making dynamic memory calls unsafe and contextually impossible.

Solution:

We must manually act as the memory manager and write bytes directly to a specific address in SRAM.

'h', 'a', 'c', 'k', ...

SRAM (Writable) at 0x20000000

The Injection (Writing to SRAM)



Since we can't write to Flash (Read-Only), we write to the **base of SRAM** (0x20000000).

```
(gdb) set {char[14]} 0x20000000 = {'h','a','c','k','y',' ',' ',' ',' ','w','o','r','l','d','\r','\0'}
```

The Result:



Read-Only

0x100019cc

Flash (0x10...)

"hello, world\r\0"



Still holds **"hello, world"**
(Untouched).



Read-Write

0x20000000

SRAM (0x20...)

"hacky, world\r\0"



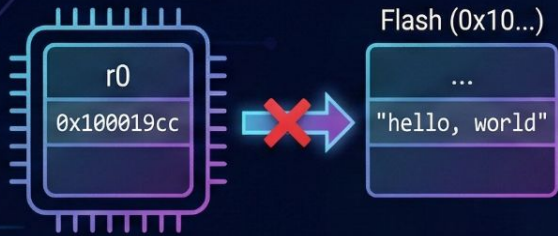
Now holds **"hacky, world"**
(Our injected payload).



The Hijack (Register Manipulation)

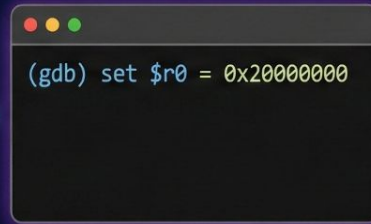
Now that our data is in memory, we must tell the CPU to use it.

Step 1: Paused at Breakpoint



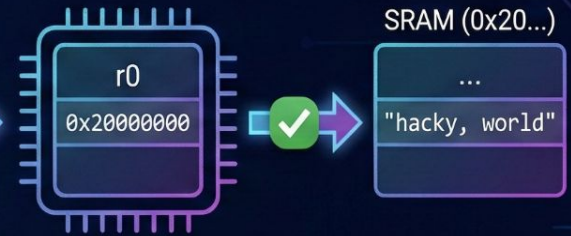
r0 currently points to 0x100019cc (Flash).

Step 2: Execute GDB Command



We execute `set $r0 = 0x20000000`.

Step 3: Register Updated



r0 now points to the start of SRAM.

Execution Flow:



puts()



When we type `continue`, the `puts` function looks at `r0`, sees the SRAM address, and prints our hacked string instead of the original.

Key Takeaways for Week 2

Hardware Constraints Matter



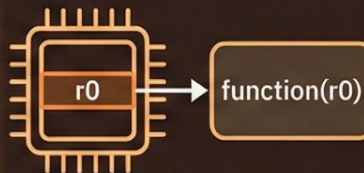
You cannot write to Flash execution memory at runtime; you must use RAM.

Bare Metal Realities



There is no OS to help you. You must manage memory and addresses manually.

Registers are King



Controlling `r0` (the function argument register) gives you control over function behavior.

Static vs. Dynamic



Ghidra
(Static)

Helps us plan the attack by finding addresses.



GDB
(Dynamic)

Allows us to execute the attack by modifying the live state.

